# KRYLOV SUBSPACE METHODS FOR CAUCHY SINGULAR INTEGRAL EQUATIONS 

P. Junghanns and K. Rost<br>Dedicated to our friend Giuseppe Mastroianni on the occasion of his 65th birthday


#### Abstract

We discuss the applicability of Krylov subspace methods to the solution of linear systems of equations occurring in the application of collocation methods to the numerical solution of Cauchy singular integral equations on an interval. The aim is to obtain fast algorithms for the approximate solution of such equations with $O(n \log n)$ computational complexity.


## 1. Introduction

During the last 15 years, the problem of the construction of fast solvers for linear systems of algebraic equations occurring in the numerical solution of singular integral and related equations have been played an important role. Since, usually, the matrices of such systems are not sparse one has to try to employ some structural and/or smoothing properties of the operators (or matrices) involved in the original equation. The classes of singular integral equations, for which there exist essential contributions in this direction can be divided into two cases, the periodic case and a case, which we call quasiperiodic.

As an example for the periodic case, let us consider a Cauchy singular integral equation on the unit circle $\mathbb{T}=\{z \in \mathbb{C}:|z|=1\}$ of the form

$$
\begin{equation*}
(A u)(t):=a(t) u(t)+\frac{b(t)}{\pi \mathrm{i}} \int_{\mathbb{T}} \frac{u(s) d s}{s-t}=f(t), \quad t \in \mathbb{T}, \tag{1.1}
\end{equation*}
$$

[^0]where $a, b, f: \mathbb{T} \longrightarrow \mathbb{C}$ are given, say, continuous functions satisfying $[a(t)]^{2}-$ $[b(t)]^{2} \neq 0, t \in \mathbb{T}$. A collocation method for equation (1.1) can be described in the following way. Let $t_{n j}=e^{2 \pi \mathrm{i} j /(2 n+1)}, j=-n, \ldots, n$, be the collocation points, $e_{j}(t)=t^{j}$, and
$$
P_{n}^{\mathbb{T}} u=\sum_{j=-n}^{n}\left\langle u, e_{j}\right\rangle_{\mathbb{T}} e_{j}
$$
the respective orthoprojection from $\boldsymbol{L}^{2}(\mathbb{T})$ onto span $\left\{e_{-n}, \ldots, e_{n}\right\}$, where
$$
\langle u, v\rangle_{\mathbb{T}}=\frac{1}{2 \pi} \int_{-\pi}^{\pi} u\left(e^{\mathrm{i} s}\right) \overline{v\left(e^{\mathrm{i} s}\right)} d s
$$

Look for an approximate solution $u_{n} \in \operatorname{im} P_{n}^{\mathbb{T}}$, such that

$$
\begin{equation*}
\left(A u_{n}\right)\left(t_{n j}\right)=f\left(t_{n j}\right), \quad j=-n, \ldots, n . \tag{1.2}
\end{equation*}
$$

If we define the interpolation operator $L_{n}^{\mathbb{T}}$ by

$$
L_{n}^{\mathbb{T}} f \in \operatorname{im} P_{n}^{\mathbb{T}}, \quad\left(L_{n}^{\mathbb{T}} f\right)\left(t_{n j}\right)=f\left(t_{n j}\right), j=-n, \ldots, n,
$$

the collocation conditions (1.2) can be written equivalently as the operator equation

$$
L_{n}^{\mathbb{T}} A u_{n}=L_{n}^{\mathbb{T}} f, \quad u_{n} \in \operatorname{im} P_{n}^{\mathbb{T}} .
$$

Write the operator $A$ of (1.1) in the form $A=a I+b S_{\mathbb{T}}$. It is well known that $S_{\mathbb{T}}: \boldsymbol{L}^{2}(\mathbb{T}) \longrightarrow \boldsymbol{L}^{2}(\mathbb{T})$ is continuous with $S_{\mathbb{T}}^{2}=I$, such that

$$
P_{\mathbb{T}}=\frac{1}{2}\left(I+S_{\mathbb{T}}\right) \quad \text { and } \quad Q_{\mathbb{T}}=\frac{1}{2}\left(I-S_{\mathbb{T}}\right)=I-P_{\mathbb{T}}
$$

are projections and $A=c P_{\mathbb{T}}+d Q_{\mathbb{T}}$, where $c=a+b, d=a-b$. Hence the collocation method (1.2) is equivalent to

$$
\begin{equation*}
L_{n}^{\mathbb{T}}\left(c P_{\mathbb{T}}+d Q_{\mathbb{T}}\right) u_{n}=L_{n}^{\mathbb{T}} f, \quad u_{n} \in \operatorname{im} P_{n}^{\mathbb{T}} . \tag{1.3}
\end{equation*}
$$

The application of Amosov's idea (see $[1,5]$ ) to this situation leads to the following method.

Choose a positive integer $m$ such that $m<n$ and $(2 n+1) /(2 m+1)$ is an integer, and look for an approximate solution $v_{n}$ as $v_{n}=P_{m}^{\mathbb{T}} w_{m}+Q_{m}^{\mathbb{T}} z_{n}$, where $Q_{m}^{\mathbb{T}}=I-P_{m}^{\mathbb{T}}, w_{m} \in \operatorname{im} P_{m}^{\mathbb{T}}, z_{n} \in \operatorname{im} P_{n}^{\mathbb{T}}$, and

$$
z_{n}=L_{n}^{\mathbb{T}} B L_{n}^{\mathbb{T}} f, \quad B=c^{-1} P_{\mathbb{T}}+d^{-1} Q_{\mathbb{T}},
$$

$$
L_{m}^{\mathbb{T}}\left(c P_{\mathbb{T}}+d Q_{\mathbb{T}}\right) w_{m}=L_{m}^{\mathbb{T}} g, \quad g=f-\left(c P_{\mathbb{T}}+d Q_{\mathbb{T}}\right) Q_{m}^{\mathbb{T}} z_{n}
$$

Using the Fourier matrix $\boldsymbol{F}_{n}=\left[t_{n j}^{k}\right]_{j, k=-n}^{n}$ of order $2 n+1$ together with its inverse $\boldsymbol{F}_{n}^{-1}=\frac{1}{2 n+1}\left[t_{n j}^{-k}\right]_{j, k=-n}^{n}$ the method can be realized with $O(n \log n)$ complexity if $m$ is chosen in such a way that $m \sim n^{1 / 3}$ (see the following algorithm). Let $\widetilde{f}_{n}=\left[f\left(t_{n j}\right)\right]_{j=-n}^{n}$ and $\widehat{f}_{n}=\left[\left\langle f, e_{j}\right\rangle_{\mathbb{T}}\right]_{j=-n}^{n}$.

## Algorithm:

1. Compute $\widehat{z}_{n}=\boldsymbol{F}_{n}^{-1}\left(\boldsymbol{C}_{n}^{-1} \boldsymbol{F}_{n} \boldsymbol{I}_{n}^{+}+\boldsymbol{D}_{n}^{-1} \boldsymbol{F}_{n} \boldsymbol{I}_{n}^{-}\right) \boldsymbol{F}_{n}^{-1} \widetilde{f}_{n}$, where $\boldsymbol{C}_{n}=$ $\operatorname{diag} \widetilde{c}_{n}, \boldsymbol{D}_{n}=\operatorname{diag} \widetilde{d}_{n}$,

$$
\boldsymbol{I}_{n}^{+}=\left[\begin{array}{cc}
\mathbf{0} & \mathbf{0} \\
\mathbf{0} & \boldsymbol{I}_{n+1}
\end{array}\right], \quad \boldsymbol{I}_{n}^{-}=\left[\begin{array}{cc}
\boldsymbol{I}_{n} & \mathbf{0} \\
\mathbf{0} & \mathbf{0}
\end{array}\right]
$$

2. Compute $\widetilde{g}_{m}=\widetilde{f}_{m}-\left(\boldsymbol{C}_{m} \boldsymbol{I}_{m, n} \boldsymbol{F}_{n} \boldsymbol{I}_{n}^{+}+\boldsymbol{D}_{m} \boldsymbol{I}_{m, n} \boldsymbol{F}_{n} \boldsymbol{I}_{n}^{-}\right) \boldsymbol{I}_{n}^{m} \widehat{z}_{n}$, where

$$
\boldsymbol{I}_{n}^{m}=\operatorname{diag}[\underbrace{1, \ldots, 1}_{n-m}, 0, \ldots, 0, \underbrace{1, \ldots, 1}_{n-m}], \quad\left[\boldsymbol{I}_{m, n}\right]_{j k}=\left\{\begin{array}{lll}
1 & : & t_{m j}=t_{n k} \\
0 & : & \text { otherwise }
\end{array}\right.
$$

3. Solve $L_{m}^{\mathbb{T}}\left(c P_{\mathbb{T}}+d Q_{\mathbb{T}}\right) w_{m}=L_{m}^{\mathbb{T}} g$.

In case of stability of the collocation method (1.3) and of appropriate smoothness properties of $a, b, f$ the solution $v_{n}$ of the above fast algorithm satisfies the same convergence rate as the solution $u_{n}$ of the usual collocation method (1.3). A consequence of this is that this fast algorithm is restricted to the case of smooth coefficients $a, b$. For more details and further fast algorithms we refer the reader to [5] and two the overview paper [14] as well as to $[4,18]$.

The Cauchy singular integral operator $S_{\mathbb{T}}$ on the unit circle satisfies the nice relations

$$
\left(S_{\mathbb{T}} e_{j}\right)(t)=\frac{1}{\pi \mathrm{i}} \int_{\mathbb{T}} \frac{e_{j}(s) d s}{s-t}=\left\{\begin{array}{cl}
e_{j}(t) & : j=0,1,2, \ldots \\
-e_{j}(t) & : j=-1,-2, \ldots
\end{array}\right.
$$

Cases, in which analogous (trigonometric) relations are available, we call quasi-periodic. One of such relations, for example, is the relation between the normed Chebyshev polynomials $W_{n}(x)$ and $V_{n}(x)$ of degree $n$ with respect
to the weights $(1-x)^{-1 / 2}(1+x)^{1 / 2}$ and $(1-x)^{1 / 2}(1+x)^{-1 / 2}$, respectively, namely

$$
\frac{1}{\pi} \int_{-1}^{1} \frac{W_{n}(y)}{y-x} \sqrt{\frac{1+y}{1-y}} d y=V_{n}(x), \quad-1<x<1, n=0,1,2, \ldots
$$

Based on this relation on can design an analogous fast algorithm as in the periodic case to solve the equation

$$
\frac{1}{\pi} \int_{-1}^{1}\left[\frac{1}{y-x}+h(x, y)\right] u(y) \sqrt{\frac{1+y}{1-y}} d y=f(x), \quad-1<x<1
$$

if the kernel $h(x, y)$ and the right hand side $f(x)$ are sufficiently smooth. In such an algorithm there are used fast trigonometric transforms instead of the discrete Fourier transform. For the interested reader we refer to [2, 9, 14]. Analogous ideas are applicable, for example, to singular integro-differential equations of Prandtl's type (see [6]).

In the present paper we discuss the applicability of Krylov subspace methods to solve systems of linear algebraic equations generated by collocation methods applied to equations of the form (2.1). Here only the structure of the matrix is used and no smoothness properties on the data are necessary. Algorithms with $O\left(n^{2}\right)$ complexity were already used (see, for example, $[3,8,10])$. In this paper our aim is to design algorithms of $O(n \log n)$ complexity.

The paper is organized as follows. In the next section we discuss the stability of the collocation methods under consideration together with the condition of the respective sequences of systems of linear equations. In Section 3 the structure of the system matrices is described and representations of these matrices are presented, which involve only fast trigonometric transforms and diagonal matrices. In Section 4 we shortly describe the Krylov subspace methods, for which we present numerical examples in Section 5.

## 2. The Collocation Method and Its Stability

For given piecewise continuous functions $a, b:[-1,1] \longrightarrow \mathbb{C}$, let us consider the Cauchy singular integral equation

$$
\begin{equation*}
a(x) u(x)+\frac{b(x)}{\pi \mathrm{i}} \int_{-1}^{1} \frac{u(y) d y}{y-x}=f(x), \quad-1<x<1 \tag{2.1}
\end{equation*}
$$

where also $f:[-1,1] \longrightarrow \mathbb{C}$ is a given function belonging to the Hilbert space $\boldsymbol{L}_{\sigma}^{2}$ of square integrable with respect to the Chebyshev weight $\sigma(x)=$
$1 / \sqrt{1-x^{2}}$ functions on $[-1,1]$. The inner product in $\boldsymbol{L}_{\sigma}^{2}$ is defined by

$$
\langle u, v\rangle_{\sigma}=\int_{-1}^{1} u(x) \overline{v(x)} \sigma(x) d x .
$$

Using the operators $a I$ and $b I$ of multiplication by $a(x)$ and $b(x)$, respectively, as well as the Cauchy singular integral operator $S$ on the interval $[-1,1]$ equation (2.1) can be written as operator equation

$$
a u+b S u=f .
$$

The operators of multiplication by a piecewise continuous function as well as the singular integral operator $S$ are linear bounded operators in $\boldsymbol{L}_{\sigma}^{2}$.

Let us solve equation (2.1) numerically. Firstly we look for an approximate solution $u_{n}(x)$ as weighted polynomial

$$
\begin{equation*}
u_{n}(x)=\sum_{k=1}^{n} \xi_{n k} \widetilde{\ell}_{n k}^{\omega}(x)=\sum_{j=0}^{n-1} a_{n j} \widetilde{u}_{j}(x), \tag{2.2}
\end{equation*}
$$

where

$$
\tilde{\ell}_{n k}^{\omega}(x)=\frac{\varphi(x) \ell_{n k}^{\omega}(x)}{\varphi\left(x_{n k}^{\omega}\right)}, \quad \widetilde{u}_{j}(x)=\varphi(x) U_{j}(x), \quad \varphi(x)=\sqrt{1-x^{2}},
$$

and $\ell_{n k}^{\omega}(x), k=1, \ldots, n$, denote the usual fundamental Lagrange interpolation polynomials with respect to the nodes $x_{n k}^{\omega}$ (the zeros of the $n$th orthogonal polynomial with respect to the weight $\omega(x)$ ). By $U_{j}(x)$ we denote the normalized Chebyshev polynomial of second kind and degree $j$, i.e.,

$$
U_{j}(\cos s)=\sqrt{\frac{2}{\pi}} \frac{\sin (j+1) s}{\sin s} .
$$

For $\omega$ we choose $\omega=\sigma$ or $\omega=\varphi$. Remark that

$$
x_{n k}^{\sigma}=\cos \frac{2 k-1}{2 n} \pi \quad \text { and } \quad x_{n k}^{\varphi}=\cos \frac{k \pi}{n+1} .
$$

Secondly we apply a collocation method

$$
\begin{equation*}
a\left(x_{n j}^{\omega}\right) u_{n}\left(x_{n j}^{\omega}\right)+\frac{b\left(x_{n j}^{\omega}\right)}{\pi \mathrm{i}} \int_{-1}^{1} \frac{u_{n}(y) d y}{y-x_{n j}^{\omega}}=f\left(x_{n j}^{\omega}\right), \quad j=1, \ldots, n, \tag{2.3}
\end{equation*}
$$

to determine the unknown coefficients $\xi_{n k}, k=1, \ldots, n$, or $a_{n j}, j=0,1, \ldots$, $n-1$, in (2.2). If we define the weighted interpolation operator $M_{n}^{\omega}=$
$\varphi L_{n}^{\omega} \varphi^{-1} I$, where $L_{n}^{\omega}$ is the usual Lagrange interpolation operator with respect to the nodes $x_{n k}^{\omega}, k=1, \ldots, n$, the collocation method (2.3) can be written as operator equation $A_{n}^{\omega} u_{n}=M_{n}^{\omega} f$, where $A_{n}^{\omega}=M_{n}^{\omega} A L_{n}$ and $L_{n}: \boldsymbol{L}_{\sigma}^{2} \longrightarrow \boldsymbol{L}_{\sigma}^{2}$ is the projection

$$
L_{n} u=\sum_{j=0}^{n-1}\left\langle u, \widetilde{u}_{j}\right\rangle_{\sigma} \widetilde{u}_{j} .
$$

Remark that the system $\left(\widetilde{u}_{n}\right)_{n=0}^{+\infty}$ forms an orthonormal basis in $\boldsymbol{L}_{\sigma}^{2}$. We call a sequence $\left(A_{n}\right)$ of operators $A_{n}: \operatorname{im} L_{n} \longrightarrow \operatorname{im} L_{n}$ stable in $\boldsymbol{L}_{\sigma}^{2}$ if, for all sufficiently large $n$, these operators are invertible and the norms of their inverses are uniformly bounded.

Theorem 2.1. ([13], Thm. 2.1) For piecewise continuous functions $a, b$ : $[-1,1] \longrightarrow \mathbb{C}$, the sequence $\left(M_{n}^{\omega}(a I+b S) L_{n}\right)$ is stable in $\boldsymbol{L}_{\sigma}^{2}$ if and only if,
(a) in the case $\omega=\sigma$, the operator $a I+b S: \boldsymbol{L}_{\sigma}^{2} \longrightarrow \boldsymbol{L}_{\sigma}^{2}$ is invertible,
(b) in the case $\omega=\varphi$, the operators $a I \pm b S: \boldsymbol{L}_{\sigma}^{2} \longrightarrow \boldsymbol{L}_{\sigma}^{2}$ are both invertible.

We remark that the collocation methods under consideration in [13] as well as in the present paper are special cases of the collocation methods investigated in [11] and [12].

Let $\lambda_{n k}^{\omega}, k=1, \ldots, n$, denote the Christoffel numbers with respect to the weight $\omega(x)$, i.e. $\lambda_{n k}^{\sigma}=\pi / n$ and $\lambda_{n k}^{\varphi}=\pi\left[\varphi\left(x_{n k}^{\varphi}\right)\right]^{2} /(n+1)$, and let $\psi_{n, \varphi}=\sqrt{(n+1) / \pi}$. By means of the Gaussian rule one can easily prove that the system $\left\{\psi_{n, \varphi} \widetilde{\varphi}_{n k}^{\varphi}: k=1, \ldots, n\right\}$ forms an orthonormal basis in the space $\operatorname{im} L_{n}$. The matrix representation of the operator $A_{n}^{\omega}$ with respect to this basis is equal to $\boldsymbol{A}_{n}^{\varphi}=\left[\alpha_{j k}^{\varphi}\right]_{j, k=1}^{n}$ with $\alpha_{j k}^{\varphi}=\left(A \tilde{\ell}_{n k}^{\varphi}\right)\left(x_{n j}^{\varphi}\right)$. Hence, the spectral norm $\left\|\boldsymbol{A}_{n}^{\varphi}\right\|$ of the matrix $\boldsymbol{A}_{n}^{\varphi}$ is equal to the norm $\left\|A_{n}^{\varphi}\right\|_{\boldsymbol{L}_{\sigma}^{2} \rightarrow \boldsymbol{L}_{\sigma}^{2}}$ of the operator $A_{n}^{\varphi}: \operatorname{im} L_{n} \longrightarrow \operatorname{im} L_{n}$. Let $\ell_{n}^{2}$ denote the $n$-dimensional complex vector space equipped with the norm

$$
\left\|\left(\xi_{k}\right)_{k=1}^{n}\right\|_{\ell_{n}^{2}}=\left(\sum_{k=1}^{n}\left|\xi_{k}\right|^{2}\right)^{1 / 2}
$$

As a consequence of Theorem 2.1 we get the following corollary.
Corollary 2.2. Let $a, b:[-1,1] \longrightarrow \mathbb{C}$ be piecewise continuous functions. The sequence $\left(\boldsymbol{A}_{n}^{\varphi}\right)$ of matrices considered as operators $\boldsymbol{A}_{n}^{\varphi}: \ell_{n}^{2} \longrightarrow \ell_{n}^{2}$ is stable if and only if both operators aI $\pm b S: \boldsymbol{L}_{\sigma}^{2} \longrightarrow \boldsymbol{L}_{\sigma}^{2}$ are invertible.

Let us mention that the operators $\boldsymbol{A}_{n}^{\varphi}$ (as well as the operators $\boldsymbol{A}_{n}^{\sigma}$ ) are strongly convergent (see [13]), and consequently, their norms are uniformly bounded. Thus, Corollary 2.2 implies that stability of the collocation method w.r.t. the Chebyshev nodes of second kind is necessary and sufficient for the boundedness of the sequence of the spectral condition numbers of $\boldsymbol{A}_{n}^{\varphi}$.

In order to get an analogous result in case of $\omega=\sigma$, let us denote by $\boldsymbol{U}_{n}^{\omega}=\left[\tau_{j k}\right]_{j, k=1}^{n}$ the matrix of the basis transformation (in im $L_{n}$ ) from $\left\{\widetilde{u}_{0}, \ldots, \widetilde{u}_{n-1}\right\}$ to $\left\{\psi_{\omega, n} \widetilde{\ell}_{n 1}^{\omega}, \ldots, \psi_{\omega, n} \widetilde{\ell}_{n n}^{\omega}\right\}$, where $\psi_{\sigma, n}=\sqrt{n / \pi}$. Moreover, let $T_{n}(x)$ denote the normalized Chebyshev polynomials of first kind and of degree $n$,

$$
T_{0}(x)=\sqrt{\frac{1}{\pi}}, \quad T_{n}(\cos s)=\sqrt{\frac{2}{\pi}} \cos n s, \quad n=1,2, \ldots
$$

In the case $\omega=\varphi$, we have $\tau_{j k}=\left\langle\psi_{\varphi, n} \widetilde{\ell}_{n k}^{\varphi}, \widetilde{u}_{j-1}\right\rangle_{\sigma}$. Thus,

$$
\begin{equation*}
\tau_{j k}=\sqrt{\frac{n+1}{\pi}} \frac{1}{\varphi\left(x_{n k}^{\varphi}\right)}\left\langle\ell_{n k}^{\varphi}, U_{j-1}\right\rangle_{\varphi}=\sqrt{\frac{2}{n+1}} \sin \frac{j k \pi}{n+1} \tag{2.4}
\end{equation*}
$$

i.e., $\boldsymbol{U}_{n}^{\varphi}=\sqrt{\frac{2}{n+1}} \boldsymbol{S}_{n}^{1}$, where $\boldsymbol{S}_{n}^{1}=\left[\sin \frac{j k \pi}{n+1}\right]_{j, k=1}^{n}$ is the first discrete sine transform (DST-1) of order $n$ (cf. [7]). Of course, due to the fact, that both bases are orthonormal ones, the matrix $\boldsymbol{U}_{n}^{\varphi}$ is unitary. Using the relation

$$
\left(1-x^{2}\right) U_{n-1}(x)=\frac{1}{2}\left[T_{n-1}(x)-T_{n+1}(x)\right], \quad n=2,3, \ldots
$$

in case $\omega=\sigma$ we get

$$
\begin{aligned}
\tau_{j k} & =\left\langle\psi_{\sigma, n} \widetilde{\ell}_{n k}^{\sigma}, \widetilde{u}_{j-1}\right\rangle_{\sigma}=\sqrt{\frac{n}{\pi}} \frac{1}{\varphi\left(x_{n k}^{\sigma}\right)}\left\langle\ell_{n k}^{\sigma}, \varphi^{2} U_{j-1}\right\rangle_{\sigma} \\
& = \begin{cases}\sqrt{\frac{\pi}{n}} \varphi\left(x_{n k}^{\sigma}\right) U_{j-1}\left(x_{n k}^{\sigma}\right) \quad: & j=1,2, \ldots, n-1 \\
\sqrt{\frac{\pi}{n}} \frac{1}{\varphi\left(x_{n k}^{\sigma}\right)} \frac{1}{2} T_{n-1}\left(x_{n k}^{\sigma}\right) & : \\
& =\left\{\begin{array}{ll}
\sqrt{\frac{2}{n}} \sin \frac{j(2 k-1) \pi}{2 n} & : \\
\\
\frac{1}{2} \sqrt{\frac{2}{n}} \sin \frac{n(2 k-1) \pi}{2 n} & :
\end{array} \quad j=1,2, \ldots, n-1\right.\end{cases} \\
& : \quad \begin{array}{ll}
2=n
\end{array}
\end{aligned}
$$

Hence, $\boldsymbol{U}_{n}^{\sigma}=\sqrt{\frac{2}{n}} \boldsymbol{D}_{n} \boldsymbol{S}_{n}^{2}$, where $\boldsymbol{S}_{n}^{2}=\left[\sin \frac{j(2 k-1) \pi}{2 n}\right]_{j, k=1}^{n}$ is the second discrete sine transform (DST-2) of order $n$ and $\boldsymbol{D}_{n}=\operatorname{diag}\left[1, \ldots, 1, \frac{1}{2}\right] \in$ $\mathbb{R}^{n \times n}$. Since, for $j=1, \ldots, n$,

$$
\widetilde{u}_{j-1}(x)=\frac{\psi_{\sigma, n}}{\psi_{\sigma, n}} \sum_{k=1}^{n} \widetilde{u}_{j-1}\left(x_{n k}^{\sigma}\right) \widetilde{\ell}_{n k}^{\sigma}(x)=\sqrt{\frac{2}{n}} \sum_{k=1}^{n} \sin \frac{j(2 k-1) \pi}{2 n} \psi_{\sigma, n} \widetilde{\ell}_{n k}^{\sigma}(x),
$$

the inverse matrix of $\boldsymbol{U}_{n}^{\sigma}$ is given by

$$
\left(\boldsymbol{U}_{n}^{\sigma}\right)^{-1}=\left[\sqrt{\frac{2}{n}} \sin \frac{j(2 k-1) \pi}{2 n}\right]_{k, j=1}^{n}=\sqrt{\frac{2}{n}}\left(\boldsymbol{S}_{n}^{2}\right)^{T}=\left(\boldsymbol{U}_{n}^{\sigma}\right)^{T} \boldsymbol{D}_{n}^{-1} .
$$

The matrix representation $\boldsymbol{A}_{n}^{\sigma}$ of the operator $A_{n}^{\sigma}: \operatorname{im} L_{n} \longrightarrow \operatorname{im} L_{n}$ with respect to the basis $\left\{\psi_{\sigma, n} \widetilde{\mathscr{\ell}}_{n k}^{\sigma}: k=1, \ldots, n\right\}$ is equal to $\boldsymbol{A}_{n}^{\sigma}=\left[\alpha_{j k}^{\sigma}\right]_{j, k=1}^{n}$ with $\alpha_{j k}^{\sigma}=\left(A \widetilde{\ell_{n k}^{\sigma}}\right)\left(x_{n j}^{\sigma}\right)$. If $\boldsymbol{A}_{n}$ denotes the matrix representation of $A_{n}$ : $\operatorname{im} L_{n} \longrightarrow \operatorname{im} L_{n}$ with respect to the orthonormal basis $\left\{\widetilde{u}_{0}, \ldots, \widetilde{u}_{n-1}\right\}$, then

$$
\begin{equation*}
\boldsymbol{A}_{n}^{\sigma}=\left(\boldsymbol{U}_{n}^{\sigma}\right)^{T} D_{n}^{-1} \boldsymbol{A}_{n} \boldsymbol{U}_{n}^{\boldsymbol{\sigma}} . \tag{2.6}
\end{equation*}
$$

If, for some nonnegative constants $c_{n}$ and $d_{n}$,

$$
c_{n}\left\|\xi_{n}\right\|_{\ell_{n}^{2}} \leq\left\|\boldsymbol{A}_{n} \xi_{n}\right\|_{\ell_{n}^{2}} \leq d_{n}\left\|\xi_{n}\right\|_{\ell_{n}^{2}} \quad \forall \xi_{n} \in \ell_{n}^{2},
$$

then, by means of (2.6), one can easily prove that

$$
\frac{c_{n}}{2}\left\|\xi_{n}\right\|_{\ell_{n}^{2}} \leq\left\|\boldsymbol{A}_{n}^{\sigma} \xi_{n}\right\|_{\ell_{n}^{2}} \leq 2 d_{n}\left\|\xi_{n}\right\|_{\ell_{n}^{2}} \quad \forall \xi_{n} \in \ell_{n}^{2} .
$$

Consequently, we have the following corollary of Theorem 2.1.
Corollary 2.3. Let $a, b:[-1,1] \longrightarrow \mathbb{C}$ be piecewise continuous functions. The sequence $\left(\boldsymbol{A}_{n}^{\sigma}\right)$ of matrices considered as operators $\boldsymbol{A}_{n}^{\sigma}: \ell_{n}^{2} \longrightarrow \ell_{n}^{2}$ is stable if and only if the operator $a I+b S: \boldsymbol{L}_{\sigma}^{2} \longrightarrow \boldsymbol{L}_{\sigma}^{2}$ is invertible.

## 3. The Structure of the Discretized Equations

With the help of well known relations between the Chebyshev polynomials of first and second kind and the Cauchy singular integral operator $S$ one can show that $\boldsymbol{A}_{n}^{\omega}$ can be written in the form (see [10, Section 5])

$$
\boldsymbol{A}_{n}^{\omega}=\boldsymbol{a}_{n}^{\omega}+\boldsymbol{b}_{n}^{\omega} \boldsymbol{S}_{n} \boldsymbol{D}_{n}(\beta),
$$

where

$$
\begin{gathered}
\boldsymbol{a}_{n}^{\sigma}=\operatorname{diag}\left[a\left(x_{n k}^{\omega}\right)-\frac{x_{n k}^{\sigma} b\left(x_{n k}^{\sigma}\right)}{n \varphi\left(x_{n k}^{\sigma}\right)}\right]_{k=1}^{n}, \quad \boldsymbol{a}_{n}^{\varphi}=\operatorname{diag}\left[a\left(x_{n k}^{\varphi}\right)\right]_{k=1}^{n}, \\
\boldsymbol{b}_{n}^{\omega}=\operatorname{diag}\left[b\left(x_{n k}^{\omega}\right)\right]_{k=1}^{n}, \boldsymbol{S}_{n}=\left[\frac{\alpha_{k}^{\omega}-\alpha_{j}^{\omega}}{x_{n k}^{\omega}-x_{n j}^{\omega}}\right]_{j, k=1}^{n}, \boldsymbol{D}_{n}(\beta)=\operatorname{diag}\left[\beta_{k}^{\omega}\right]_{k=1}^{n},
\end{gathered}
$$

and

$$
\alpha_{k}^{\sigma}=(-1)^{k} \varphi\left(x_{n k}^{\sigma}\right), \quad \beta_{k}^{\sigma}=\frac{(-1)^{k}}{n \mathrm{i}}, \quad \alpha_{k}^{\varphi}=(-1)^{k}, \quad \beta_{k}^{\varphi}=\frac{(-1)^{k} \varphi\left(x_{n k}^{\varphi}\right)}{(n+1) \mathrm{i}} .
$$

Since matrices of the form $\boldsymbol{S}_{n}$ can be represented with the help of fast discrete trigonometric transforms, in [10] the authors propose to apply iterative methods to solve the systems $\boldsymbol{A}_{n}^{\omega} \xi_{n}=\eta_{n}$ of linear equations effectively. This means the following. Let

$$
C_{n}^{\omega}=\left[\frac{1}{x_{n k}^{\omega}-x_{n j}^{\omega}}\right]_{j, k=1}^{n} \quad \text { (zeros on the main diagonal). }
$$

Then $\boldsymbol{S}_{n}=\boldsymbol{C}_{n}^{\omega} \boldsymbol{D}_{n}(\alpha)-\boldsymbol{D}_{n}(\alpha) \boldsymbol{C}_{n}^{\omega}$ and ([7, Theorems 5.2, 5.5])

$$
\boldsymbol{C}_{n}^{\omega}= \begin{cases}-\frac{1}{n} \boldsymbol{D}_{n}^{\sigma} \boldsymbol{C}_{n}^{3} \boldsymbol{\Gamma}_{n}^{\sigma} \boldsymbol{C}_{n}^{2} \boldsymbol{D}_{n}^{\sigma} & : \omega=\sigma, \\ -\frac{1}{n+1} \boldsymbol{D}_{n}^{\varphi} \boldsymbol{S}_{n}^{1} \boldsymbol{\Gamma}_{n}^{\varphi} \boldsymbol{S}_{n}^{1} \boldsymbol{D}_{n}^{\varphi} & : \quad \omega=\varphi,\end{cases}
$$

where

$$
\begin{gathered}
\boldsymbol{D}_{n}^{\omega}=\operatorname{diag}\left[\frac{(-1)^{k}}{\varphi\left(x_{n k}^{\omega}\right)}\right]_{k=1}^{n}, \quad \boldsymbol{\Gamma}_{n}^{\omega}=\operatorname{tridiag}\left[\begin{array}{lll}
-\boldsymbol{c}_{\omega} & \mathbf{0} & \boldsymbol{c}_{\omega}
\end{array}\right], \\
\boldsymbol{c}_{\sigma}=\frac{1}{2}\left[\begin{array}{llll}
1 & 3 & \cdots & 2 n-3
\end{array}\right], \quad \boldsymbol{c}_{\varphi}=\frac{1}{2}\left[\begin{array}{llll}
3 & 5 & \cdots & 2 n-1
\end{array}\right],
\end{gathered}
$$

and $\boldsymbol{C}_{n}^{2}$ and $\boldsymbol{C}_{n}^{3}$ denote the second and third discrete cosine transform of order $n$, respectively,

$$
\boldsymbol{C}_{n}^{2}=\left[\cos \frac{(j-1)(2 k-1) \pi}{2 n}\right]_{j, k=1}^{n}, \quad \boldsymbol{C}_{n}^{3}=\left(\boldsymbol{C}_{n}^{2}\right)^{T} .
$$

Consequently, using these representations four trigonometric transforms are needed to apply the matrix $\boldsymbol{A}_{n}^{\omega}$ to a vector.

In what follows let us find other representations of the matrices $\boldsymbol{A}_{n}^{\sigma}$ and $\boldsymbol{A}_{n}^{\varphi}$, which allow the matrix vector multiplication with the help of only two trigonometric transforms. For this, recall the well known relation

$$
S \widetilde{u}_{m}=S \varphi U_{m}=\mathrm{i} T_{m+1}, \quad m=0,1,2, \ldots
$$

and use (2.5) and (2.4) to compute

$$
\begin{aligned}
\left(S \tilde{\ell}_{n k}^{\omega}\right)\left(x_{n j}^{\omega}\right) & =\mathrm{i} \sum_{m=1}^{n}\left\langle\widetilde{\ell}_{n k}^{\omega}, \widetilde{u}_{m-1}\right\rangle_{\sigma} T_{m}\left(x_{n j}^{\omega}\right) \\
& = \begin{cases}\frac{2 \mathrm{i}}{n} \sum_{m=1}^{n-1} \cos \frac{m(2 j-1) \pi}{2 n} \sin \frac{m(2 k-1) \pi}{2 n} & : \omega=\sigma \\
\frac{2 \mathrm{i}}{n+1} \sum_{m=1}^{n} \cos \frac{m j \pi}{n+1} \sin \frac{m k \pi}{n+1} & : \omega=\varphi\end{cases}
\end{aligned}
$$

where, in the case $\omega=\sigma$, we took into account that $T_{n}\left(x_{n j}^{\sigma}\right)=0$. Hence we have

$$
\begin{equation*}
\boldsymbol{A}_{n}^{\sigma}=\widetilde{\boldsymbol{a}}_{n}^{\sigma}-\frac{2}{n \mathrm{i}} \boldsymbol{b}_{n}^{\sigma} \boldsymbol{C}_{n}^{3} \boldsymbol{V}_{n} \boldsymbol{S}_{n}^{2} \quad \text { and } \quad \boldsymbol{A}_{n}^{\varphi}=\boldsymbol{a}_{n}^{\varphi}-\frac{2}{(n+1) \mathrm{i}} \boldsymbol{b}_{n}^{\varphi} \widetilde{\boldsymbol{C}}_{n}^{1} \boldsymbol{S}_{n}^{1} \tag{3.1}
\end{equation*}
$$

with $\boldsymbol{V}_{n}=\left[\delta_{j-1, k}\right]_{j, k=1}^{n}, \widetilde{\boldsymbol{a}}_{n}^{\sigma}=\operatorname{diag}\left[a\left(x_{n k}^{\sigma}\right)\right]_{k=1}^{n}$, and a submatrix $\widetilde{\boldsymbol{C}}_{n}^{1}=$ $\left[\cos \frac{j k \pi}{n+1}\right]_{j, k=1}^{n}$ of the first discrete cosine transform

$$
C_{n+2}^{1}=\left[\cos \frac{j k \pi}{n+1}\right]_{j, k=0}^{n+1}
$$

of order $n+2$.

## 4. The Application of Krylov Subspace Methods

Given a regular matrix $\boldsymbol{A} \in \mathbb{C}^{n \times n}$ and a vector $\eta \in \mathbb{C}^{n}$ we have to solve the system

$$
\begin{equation*}
\boldsymbol{A} \xi=\eta \tag{4.1}
\end{equation*}
$$

The basic idea of the Krylov subspace methods consists in minimizing the residual $\eta-\boldsymbol{A} \xi$ for $\xi=\xi_{0}+\zeta$, where

$$
\zeta \in \mathbb{K}_{k}\left(\boldsymbol{A}, \rho_{0}\right)=\operatorname{span}\left\{\rho_{0}, \boldsymbol{A} \rho_{0}, \ldots, \boldsymbol{A}^{k-1} \rho_{0}\right\}
$$

and $\rho_{0}=\eta-\boldsymbol{A} \xi_{0}$ (see, for example, [16]). Remark that in our situation the matrices $\boldsymbol{A}_{n}^{\omega}$ are not symmetric as well as not sparse. But, due to the representations (3.1) these matrices have a nice structure which enables us to apply them to a vector of length $n$ with $O(n \log n)$ computational complexity (cf. [19]).

In the present paper we compare the numerical results (i.e., the number of iteration steps necessary to obtain a prescribed accuracy) for three Krylov subspace methods (GMRES, FOM, and CGNR) described shortly in the following:

- The GMRES (generalized minimal residual) algorithm minimizes the norm $\|\eta-\boldsymbol{A} \xi\|\left(\|\cdot\|=\|\cdot\|_{\ell_{n}^{2}}\right)$, which is equivalent to finding $\xi_{k}=\xi_{0}+\zeta_{k}$, $\zeta_{k} \in \mathbb{K}_{k}\left(\boldsymbol{A}, \rho_{0}\right)$, such that

$$
\eta-\boldsymbol{A} \xi_{k} \perp \boldsymbol{A} \mathbb{K}_{k}\left(\boldsymbol{A}, \rho_{0}\right)
$$

Using Arnoldi's procedure with the modified Gram-Schmidt algorithm to compute an orthonormal basis in $\mathbb{K}_{k}\left(\boldsymbol{A}, \rho_{0}\right)$ leads to the following algorithm, in which $\varepsilon>0$ defines the prescribed accuracy:
$1^{\circ} \xi_{0}, \quad \rho_{0}:=\eta-\boldsymbol{A} \xi_{0}, \quad \boldsymbol{v}_{1}:=\left\|\rho_{0}\right\|^{-1} \rho_{0}, \quad k:=0$
$2^{\circ}$ (Arnoldi's procedure with modified Gram-Schmidt algorithm)

$$
\begin{aligned}
& -k:=k+1, \quad \widehat{\boldsymbol{v}}:=\boldsymbol{A} \boldsymbol{v}_{k} \\
& -i=1,2, \ldots, k: \quad h_{i k}:=\left\langle\widehat{\boldsymbol{v}}, \boldsymbol{v}_{i}\right\rangle, \quad \widehat{\boldsymbol{v}}:=\widehat{\boldsymbol{v}}-h_{i k} \boldsymbol{v}_{i} \\
& -h_{k+1, k}:=\|\widehat{\boldsymbol{v}}\|, \quad \boldsymbol{v}_{k+1}:=h_{k+1, k}^{-1} \widehat{\boldsymbol{v}}
\end{aligned}
$$

$3^{\circ} \xi_{k}:=\xi_{0}+\boldsymbol{V}_{k} \zeta_{k}$, where $\zeta_{k}$ is the solution of the minimization problem

$$
\begin{equation*}
\mathcal{J}(\zeta):=\left\|\gamma \boldsymbol{e}_{1}-\widetilde{\boldsymbol{H}}_{k} \zeta\right\| \longrightarrow \min \quad\left(\zeta \in \mathbb{C}^{k}, \gamma=\left\|\rho_{0}\right\|\right) \tag{4.2}
\end{equation*}
$$

with

$$
\widetilde{\boldsymbol{H}}_{k}=\left[\begin{array}{cccc}
h_{11} & h_{12} & \cdots & h_{1 k} \\
h_{21} & h_{22} & & h_{2 k} \\
& \ddots & \ddots & \vdots \\
& & \ddots & h_{k k} \\
0 & & & h_{k+1, k}
\end{array}\right], \quad \boldsymbol{V}_{k}=\left[\begin{array}{llll}
\boldsymbol{v}_{1} & \boldsymbol{v}_{2} & \cdots & \boldsymbol{v}_{k}
\end{array}\right]
$$

$4^{\circ}$ if $\left\|\boldsymbol{A} \xi_{k}-\eta\right\|>\varepsilon\left\|\rho_{0}\right\|$ goto 2.

To solve the least squares problem (4.2) one uses plane rotations to get a QR-factorization of $\widetilde{\boldsymbol{H}}_{k}$, say $\boldsymbol{Q}_{k} \widetilde{\boldsymbol{H}}_{k}=\boldsymbol{R}_{k}, \boldsymbol{Q}_{k}=\left(\boldsymbol{Q}_{k}^{-1}\right)^{T} \in \mathbb{C}^{(k+1) \times(k+1)}$. This factorization can be updated at each iteration step and the norm of the minimal residual is computed by $\left\|\boldsymbol{A} \xi_{k}-\eta\right\|=\left(\gamma \boldsymbol{Q}_{k} \boldsymbol{e}_{1}\right)_{k+1}$ (the last component of $\gamma \boldsymbol{Q}_{k} \boldsymbol{e}_{1}$ in the $k$ th iteration step) without computing the optimal solution $\xi_{k}$ (see [17]). If $n_{\boldsymbol{A}}$ is the number of multiplications, which are necessary to apply the matrix $\boldsymbol{A}$ (of order $n$ ) to a vector of length $n$, then $m$ (= dimension of the Krylov space) steps of GMRES require $O\left(m^{2} n\right)+m n_{\boldsymbol{A}}$ multiplications (cf. [17]).

- The FOM (full orthogonalization method) determines $\xi_{k}=\xi_{0}+\zeta_{k}$, $\zeta_{k} \in \mathbb{K}_{k}\left(\boldsymbol{A}, \rho_{0}\right)$, by the condition

$$
\eta-\boldsymbol{A} \xi_{k} \perp \mathbb{K}_{k}\left(\boldsymbol{A}, \rho_{0}\right)
$$

The algorithm can be shortly described as follows:
$1^{\circ} \xi_{0}, \quad \rho_{0}:=\eta-\boldsymbol{A} \xi_{0}, \quad \boldsymbol{v}_{1}:=\left\|\rho_{0}\right\|^{-1} \rho_{0}, \quad k:=0$
$2^{\circ}$ (Arnoldi's procedure with modified Gram-Schmidt algorithm)
$-k:=k+1, \quad \widehat{\boldsymbol{v}}:=\boldsymbol{A} \boldsymbol{v}_{k}$
$-i=1,2, \ldots, k: \quad h_{i k}:=\left\langle\widehat{\boldsymbol{v}}, \boldsymbol{v}_{i}\right\rangle, \quad \widehat{\boldsymbol{v}}:=\widehat{\boldsymbol{v}}-h_{i k} \boldsymbol{v}_{i}$
$-h_{k+1, k}:=\|\widehat{\boldsymbol{v}}\|, \quad \boldsymbol{v}_{k+1}:=h_{k+1, k}^{-1} \widehat{\boldsymbol{v}}$
$3^{\circ} \xi_{k}:=\xi_{0}+\boldsymbol{V}_{k} \zeta_{k}$, where $\zeta_{k}=\gamma \boldsymbol{H}_{k}^{-1} \boldsymbol{e}_{1}, \gamma=\left\|\rho_{0}\right\|$, and

$$
\boldsymbol{H}_{k}=\left[\begin{array}{cccc}
h_{11} & h_{12} & \cdots & h_{1 k} \\
h_{21} & h_{22} & & h_{2 k} \\
& \ddots & \ddots & \vdots \\
0 & & \ddots & h_{k k}
\end{array}\right], \quad \boldsymbol{V}_{k}=\left[\begin{array}{cccc}
\boldsymbol{v}_{1} & \boldsymbol{v}_{2} & \cdots & \boldsymbol{v}_{k}
\end{array}\right]
$$

$4^{\circ}$ if $\left\|\boldsymbol{A} \xi_{k}-\eta\right\|>\varepsilon\left\|\rho_{0}\right\|$ goto 2.
The residual norm after the $k$ th step is equal to $\left\|\boldsymbol{A} \xi_{k}-\eta\right\|=h_{k+1, k}\left|\boldsymbol{e}_{k}^{T} \zeta_{k}\right|$. Here $m$ steps of the algorithm need $O\left(m^{2} n+m^{3}\right)+m n_{A}$ multiplications.

- The CGNR method (conjugate gradient algorithm for the normal equation) is the application of the CG-algorithm to the system $\boldsymbol{A}^{T} \boldsymbol{A} \xi=\eta$, which corresponds to the least squares problem of minimizing the residual $\|\eta-\boldsymbol{A} \xi\|$. Consequently, $\xi_{k}$ is determined by the condition

$$
\boldsymbol{A}^{T}\left(\eta-\boldsymbol{A} \xi_{k}\right) \perp \mathbb{K}_{k}\left(\boldsymbol{A}^{T} \boldsymbol{A}, \boldsymbol{A}^{T} \rho_{0}\right)
$$

The algorithm can be organized in the following way (cf. [16, Sect. 8.3]):

$$
\begin{aligned}
& 1^{\circ} \xi_{0}, \quad \rho_{0}:=\eta-\boldsymbol{A} \xi_{0}, \quad \zeta_{0}:=\boldsymbol{A}^{T} \rho_{0}, \quad \boldsymbol{p}_{0}:=\zeta_{0}, k:=0 \\
& 2^{\circ} \widehat{\boldsymbol{v}}:=\boldsymbol{A} \boldsymbol{p}_{k} \\
& 3^{\circ} \alpha:=\left\|\zeta_{k}\right\|^{2} /\|\widehat{\boldsymbol{v}}\|^{2} \\
& 4^{\circ} \xi_{k+1}:=\xi_{k}+\alpha \boldsymbol{p}_{k}, \quad \rho_{k+1}:=\rho_{k}-\alpha \widehat{\boldsymbol{v}}, \\
& \quad \text { if }\left\|\rho_{k+1}\right\|<\varepsilon\left\|\rho_{0}\right\| \text { then stop } \\
& 5^{\circ} \zeta_{k+1}:=\boldsymbol{A}^{T} \rho_{k+1} \\
& 6^{\circ} \beta:=\left\|\zeta_{k+1}\right\|^{2} /\left\|\zeta_{k}\right\|^{2}, \quad \boldsymbol{p}_{k+1}:=\zeta_{k+1}+\beta \boldsymbol{p}_{k} \\
& 7^{\circ} k:=k+1 \text { goto } 2 .
\end{aligned}
$$

For $m$ steps of CGNR one needs $O(m n)+m\left(n_{\boldsymbol{A}}+n_{\boldsymbol{A}^{T}}\right)$ multiplications. If we denote by $\kappa=\operatorname{cond}(\boldsymbol{A})$ the spectral condition number of the matrix $\boldsymbol{A}$, i.e., cond $(\boldsymbol{A})=\|\boldsymbol{A}\|\left\|\boldsymbol{A}^{-1}\right\|$, where $\|\boldsymbol{A}\|$ denotes the norm of $\boldsymbol{A}: \ell_{n}^{2} \longrightarrow \ell_{n}^{2}$, then the $\ell_{n}^{2}$-norm of the residual $\rho_{m}$ after $m$ steps of CGNR can be estimated by

$$
\left\|\rho_{m}\right\| \leq 2\left(\frac{\kappa-1}{\kappa+1}\right)^{m}\left\|\rho_{0}\right\|
$$

(see [15, Section 2]). Thus, taking into account the representations (3.1), in the case of stable sequences $\left(\boldsymbol{A}_{n}^{\omega}\right)$ (cf. Corollaries 2.2 and 2.3 ), we get that the complexity of CGNR can be estimated by $O((\log \varepsilon)(n \log n))$ if $\varepsilon>0$ is the prescribed accuracy.

## 5. Numerical Examples

Let us consider three examples of Cauchy singular integral equations of the form (2.1) together with the collocation methods (2.3) and the Krylov subspace methods for the solution of the respective systems of linear equations considered in the previous section. In the first example both operators $a I+b S$ and $a I-b S$ are invertible in $\boldsymbol{L}_{\sigma}^{2}$, while in the second and third example only the operator $a I+b S$ is invertible. Due to the Corollaries 2.2 and 2.3 this means that in the last two examples the sequence of the condition numbers of the respective matrices are bounded only for collocation with respect to the Chebyshev nodes of first kind. In each example we choose $\varepsilon$ equal to $10^{-10}$. In the tables we collect the number of iterations (equal to the dimension of the Krylov space in the last step) necessary to achieve this accuracy by the different methods. For all computations the initial guess is equal to the vector $\xi_{0}=\left[\begin{array}{lll}1 & \cdots & 1\end{array}\right]^{T}$.

Example 5.1. Let us consider an example with the piecewise continuous coefficients

$$
a(x)=\left\{\begin{array}{ccc}
2 & : & x \leq-0.5 \\
3+x & : & x>-0.5
\end{array}\right\}, \quad b(x)=\mathrm{i}\left\{\begin{array}{cll}
x & : & x \leq 0.5 \\
x-1 & : & x>0.5
\end{array}\right\}
$$

and the right hand side

$$
f(x)=\left\{\begin{array}{ccc}
1 & : & x \leq 0.5 \\
x^{2}-1 & : & x>0.5
\end{array}\right\} .
$$

Collocation w.r.t. $x_{n j}^{\sigma}, k=1, \ldots, n$

|  | number of iterations |  |  |
| ---: | :---: | :---: | :---: |
| $n$ | GMRES | FOM | CGNR |
| 512 | 24 | 24 | 26 |
| 1024 | 25 | 25 | 27 |
| 16384 | 26 | 26 | 28 |
| 32768 | 26 | 26 | 29 |
| 65536 | 26 | 26 | 30 |
| 131072 | 27 | 27 | 30 |

Collocation w.r.t. $x_{n j}^{\varphi}, k=1, \ldots, n$

|  | number of iterations |  |  |
| ---: | :---: | :---: | :---: |
| $n$ | GMRES | FOM | CGNR |
| 511 | 21 | 22 | 23 |
| 1023 | 22 | 22 | 23 |
| 16383 | 22 | 22 | 24 |
| 32767 | 22 | 22 | 24 |
| 65535 | 23 | 23 | 25 |
| 131071 | 23 | 23 | 25 |

We see that the number of iterations does not essentially increase with the size $n$ of the matrices.

Example 5.2. In this example we have $a(x)=\sqrt{1-x}, b(x)=-\mathrm{i} x$, and $f(x)=|x|$.

Collocation w.r.t. $x_{n j}^{\sigma}, k=1, \ldots, n$

|  | number of iterations |  |  |
| ---: | :---: | :---: | :---: |
| $n$ | GMRES | FOM | CGNR |
| 512 | 73 | 74 | 26 |
| 1024 | 73 | 75 | 26 |
| 16384 | 74 | 76 | 27 |
| 32768 | 74 | 76 | 27 |
| 65536 | 75 | 76 | 27 |
| 131072 | 75 | 76 | 28 |

Collocation w.r.t. $x_{n j}^{\varphi}, k=1, \ldots, n$

|  | number of iterations |  |  |
| ---: | :---: | :---: | :---: |
| $n$ | GMRES | FOM | CGNR |
| 511 | 70 | 71 | 32 |
| 1023 | 70 | 71 | 32 |
| 16383 | 69 | 71 | 36 |
| 32767 | 69 | 71 | 37 |
| 65535 | 69 | 70 | 38 |
| 131071 | 69 | 70 | 39 |

Example 5.3. Let $a(x)=\sqrt{1.01-x^{2}}, b(x)=-\mathrm{i} x$, and $f(x)=|x|$.

Collocation w.r.t. $x_{n j}^{\sigma}, k=1, \ldots, n$

|  | number of iterations |  |  |
| ---: | :---: | :---: | :---: |
| $n$ | GMRES | FOM | CGNR |
| 512 | 56 | 57 | 8 |
| 1024 | 56 | 57 | 8 |
| 16384 | 56 | 57 | 9 |
| 32768 | 56 | 57 | 9 |
| 65536 | 56 | 57 | 9 |
| 131072 | 56 | 57 | 9 |

Collocation w.r.t. $x_{n j}^{\varphi}, k=1, \ldots, n$

|  | number of iterations |  |  |
| ---: | :---: | :---: | :---: |
| $n$ | GMRES | FOM | CGNR |
| 511 | 53 | 54 | 9 |
| 1023 | 53 | 53 | 9 |
| 16383 | 51 | 52 | 12 |
| 32767 | 51 | 52 | 12 |
| 65535 | 50 | 51 | 12 |
| 131071 | 50 | 51 | 13 |

We see that also in the case of increasing condition numbers the number of iterations behaves well. For this, one should take into consideration that, in case of the collocation points $x_{n j}^{\varphi}$, only one (Example 5.2) or two (Example 5.3) singular values of the matrices $\boldsymbol{A}_{n}^{\varphi}$ tend to zero and the others stay away from zero (see [13, Section 5]).

## REFERENCES

1. B. A. Amosov: On the approximate solution of elliptic pseudodifferential equations on smooth curves. Z. Anal. Anw. 9 (1990), 545-563 (in Russian).
2. D. Berthold, W. Hoppe and B. Silbermann: A fast algorithm for solving the generalized airfoil equation. J. Comput. Appl. Math. 43 (1992), 185-219.
3. D. Berthold and P. Junghanns: Direct multiple grid methods for solving singular integral equations. Wiss. Z. d. TUK, Heft 2 (1987), 180-186.
4. D. Berthold and B. Silbermann: Corrected collocation methods for periodic pseudodifferential equations. Numer. Math. 70 (1995), 397-425.
5. D. Berthold and B. Silbermann: The fast solution of periodic pseudodifferential equations. Appl. Anal. 63 (1996), 3-23.
6. M. R. Capobianco, G. Criscuolo and P. Junghanns: A fast algorithm for Prandtl's integro-differential equation. J. Comput. Appl. Math. 77 (1997), 103-128.
7. G. Heinig and K. Rost: Representations of Cauchy matrices with Chebyshev nodes using trigonometric transforms. In: Structured Matrices, Recent Advances and Applications (D. A. Bini, E. Tyrtyshnikov and P. Yalamov, eds.), Advances in the Theory of Computational Mathematics, Vol. 4, Nova Science Publishers, 2001, pp. 135-147.
8. P. Junghanns: Numerical solution of a free surface seepage problem from nonlinear channel. Appl. Anal. 63 (1996), 87-110.
9. P. Junghanns and U. Luther: Uniform convergence of a fast algorithm for Cauchy singular integral equations. Linear Algebra Appl. 275276 (1998), 327-347.
10. P. Junghanns, K. Müller and K. Rost: On collocation methods for nonlinear Cauchy singular integral equations. In: Toeplitz Matrices and Singular Integral Equations, The Bernd Silbermann Anniversary Volume (A. Böttcher, I. Gohberg and P. Junghanns,eds.), Operator Theory, Advances and Applications, Vol. 135, Birkhäuser, 2002, pp. 209-233.
11. P. Junghanns and A. Rathsfeld: On polynomial collocation for Cauchy singular integral equations with fixed singularities. Integral Equations Operator Theory 43 (2002), 155-176.
12. P. Junghanns and A. Rogozhin: Collocation methods for singular integral equations on the interval. Electr. Trans. Numer. Anal. 17 (2004), 11-75.
13. P. Junghanns, S. Roch and B. Silbermann: Collocation methods for systems of Cauchy singular integral equations on an interval. Comp. Technologies 6 (2001), 88-124.
14. P. Junghanns and B. Silbermann: Numerical analysis for onedimensional Cauchy singular integral equations. J. Comput. Appl. Math. 125 (2000), 395-421.
15. N. M. Nachtigal, S. C. Reddy and L. N. Trefethen: How fast are nonsymmetric matrix iterations?. SIAM J. Matrix Anal. Appl. 13 (1992), 778-795.
16. Y. SaAD: Iterative Methods for Sparse Linear Systems. Second edition, Society for Industrial and Applied Mathematics, Philadelphia, PA, 2003.
17. Y. SaAd and M. H. Schultz: GMRES: A generalized minimal residual algorithm for solving nonsymmetric linear systems. SIAM J. Sci. Stat. Comput. 7 (1986), 856-869.
18. J. Saranen and G. Vainikko: Fast solvers of integral and pseudodifferential equations on closed curves. Math. Comp. 67 (1998), 1473-1491.
19. G. Steidl and M. Tasche: A polynomial approach to fast algorithms for discrete Fourier-cosine and Fourier-sine transforms. Math. Comp. 56 (1991), 281-296.

Technical University of Chemnitz
Faculty of Mathematics
Department of Mathematics
D-09107 Chemnitz, Germany
e-mails: peter.junghanns@mathematik.tu-chemnitz.de
karla.rost@mathematik.tu-chemnitz.de


[^0]:    Received November 5, 2004.
    2000 Mathematics Subject Classification. Primary 65F10, 65R20; Secondary 65N12, 65 N 22 .

